

14<sup>th</sup> ANNUAL WORKSHOP 2018

### NVMf based Integration of Non-volatile Memory in a Distributed System - Lessons learned

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### Outline

- Target Distributed System: Apache Crail
- NVMf
  - Basic operation
  - Available implementations
- The Crail NVMf Tier
  - Integration with Crail
  - Implementation details
  - A new NVMf host library
- Measurements
  - Microbenchmarks
  - TeraSort
  - TPC-DS workload
- Summary & Outlook

# user level and asynchronous I/O Implements efficient distributed store for ephemeral

#### Unifies isolated, incompatible efforts to integrate high performance I/O with BigData processing frameworks

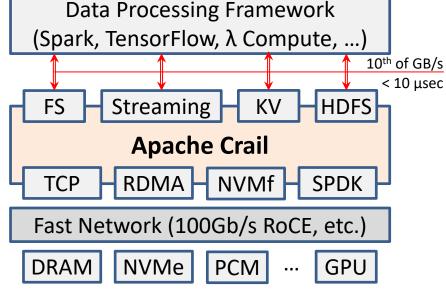
Targets I/O acceleration for Big Data frameworks due to

 Enables dramatically faster job completion

data at native hardware speeds

- Applicable to Spark, Flink, TensorFlow, Caffee, ...
- Flexibly makes best use of available I/O infrastructure
- Apache Incubator Project (since 11/17)







The Target System: Apache Crail (Incubating)

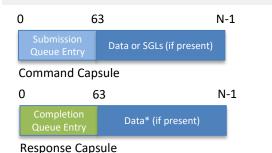
### NVMf in a Nutshell

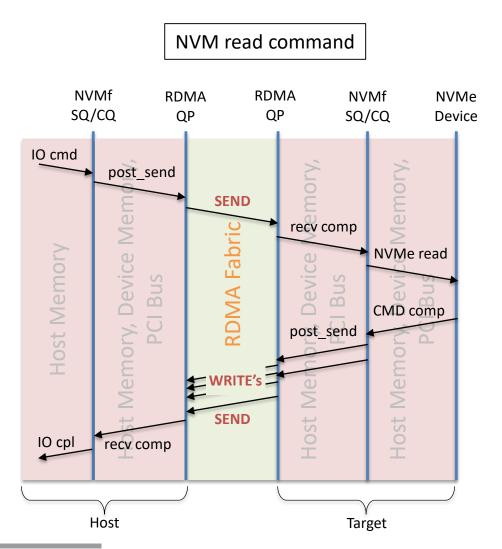
#### Extends low latency NVMe interface over distance

- Retains BW and latency (< 5us additional delay)
- Allows to retain user level I/O capability
- Message based NVMe operations
- Host/Target model
- Defined transports: RDMA (IB, RoCE, iWarp), or FC
  - Mapping of I/O Submission + Completion Queue to RDMA QP model
- Fabrics + Admin Commands
  - Discover, Connect, Authenticate, ...

#### I/O Commands

- SEND/RECV to transfer Command/Response Capsules (+ optionally data)
- READ/WRITE to transfer data

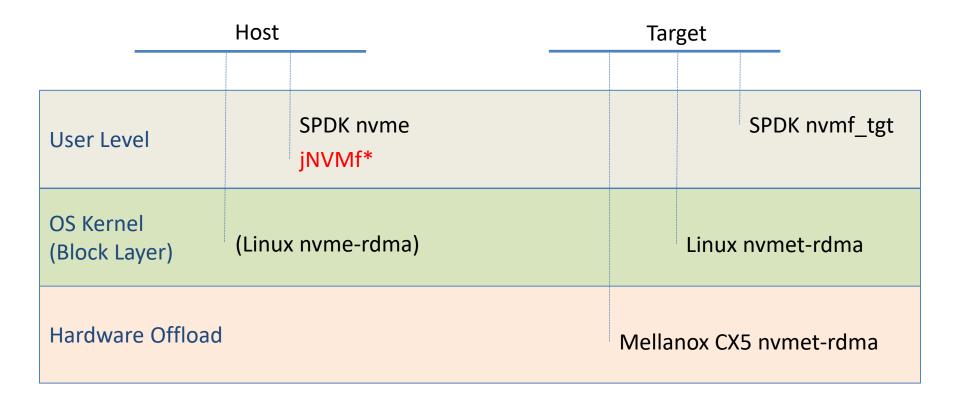




\*Not with RDMA transport

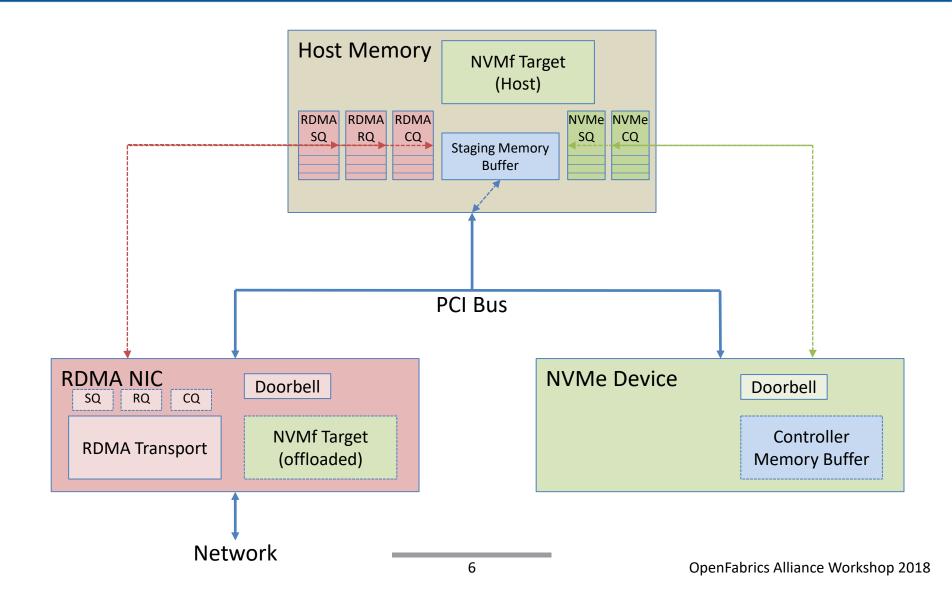
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### **NVMf Implementations available**

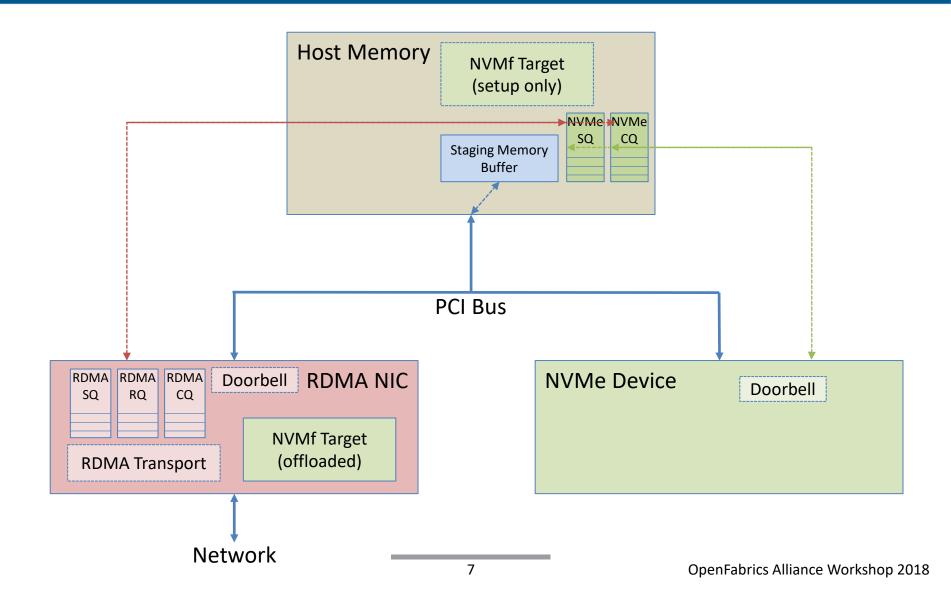


\*https://github.com/zrlio/jNVMf

### Landscape of NVMf Implementations (Target)



### Mellanox Target Offloading (w/o CMB)



### Crail: The NVMf Tier

#### Another fast storage tier in Crail

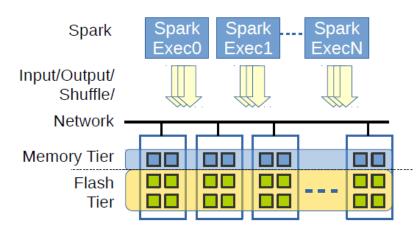
- I/O to distributed NVMe at its native speed
- Uses NMVf for remote NVM access
- User level I/O at client side

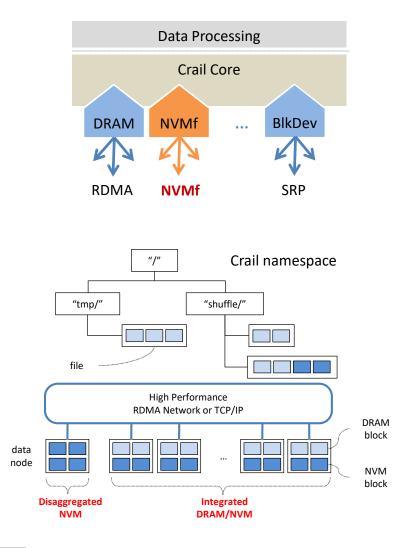
### Flexibility

- Storage disaggregation or integration
- Spill over from DRAM, or dedicated tier

#### NVMf integration

- Client library
- Data node implementation





## Putting Things together – NVMf in Crail

#### Crail Storage Tier Control

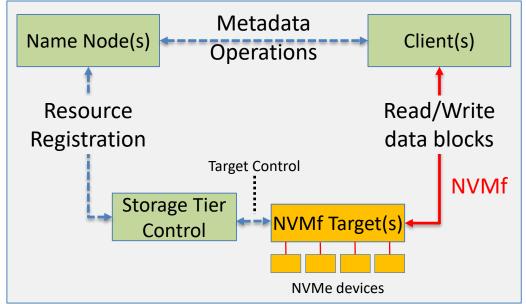
- Connects to any NVMf Target code
  - SPDK, Linux kernel, Offload
- Registers targets NVMe blocks with NameNode(s)
- Not in the block read/write NVMf fast path client ←> target

#### Crail NVMf Client

- Deploys NVMf Host code
  - SPDK host library, or
  - jNVMf host library
- Connects with NVMf Target(s)
- Gains NVMe block access info via NameNode
- Allows for application sub-block read/write access
  - Must implement RMW

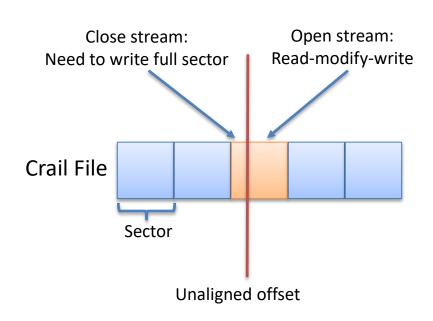
#### Crail Name Node

- Maintain storage resources
  - Get blocks registered
  - Give out blocks and back
  - May write logs



### **Crail NVMf Tier Design Discussion**

- NVMf tier must follow semantics of all Crail storage tiers:
  - Provide bytegranular access to data store, but
  - Must cope with block only media access
  - RMW implementation
    - NVMf does not support byte offset access for RDMA transport (no bit bucket type)
    - Becomes client activity above NVMf host
    - Adds another full network round trip
- Try to avoid unaligned accesses
  - Append only
  - Use buffered streams which internally aligns to block and buffer sizes
  - But cannot avoid unaligned access when closing and reopening buffered stream



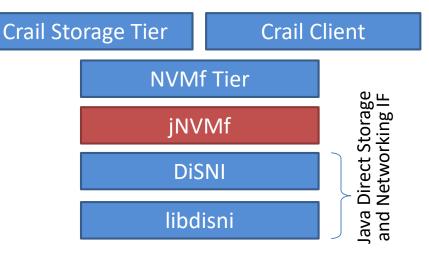
### jNVMf: Yet another NVMf Host Library?

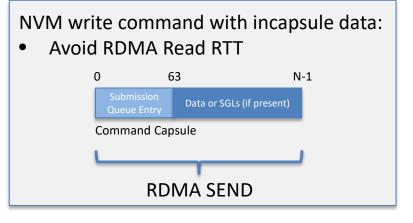
#### Dependencies and Stability

- SPDK → DPDK
- Memory management
- Bugs
- DPDK not meant to be used as a library:
  - "owns" application
  - No shared library build

#### Clean slate approach: jNVMf

- Simpler memory management
- Incapsule data support (write accel.)
- Dataset management support
- RDMA inline SEND support (read accel.)
- Reduced JNI overhead
- Easier to integrate with Crail





### **NVMe Devices we deployed**

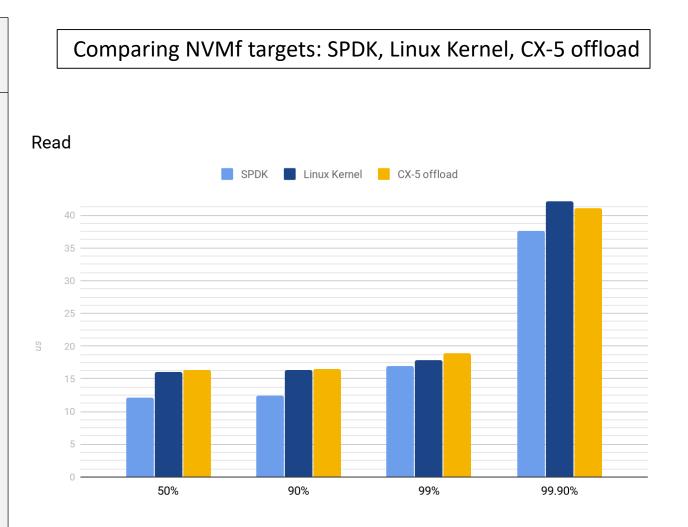
	NVMe Version	Read Latency (4K)	Write Latency (4K)	Read IOPS (4K)	Write IOPS (4K)	Read Throughput	Write Throughput
Samsung 960Pro	1.2	53.3us	7.8us	483K	386K	3186MB/s	1980MB/s
Intel 3D XPoint 900p	1.2	6.3us	7.2us	580K	557K	2586MB/s	2195MB/s

#### Many advanced NVMe features not supported:

- Scatter-gather list
- Arbitration mechanisms
- Namespace management (or only 1 namespace is supported)
- Controller memory buffer

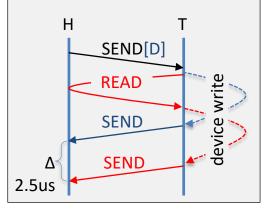
### NVMf Latency 4k (Read)

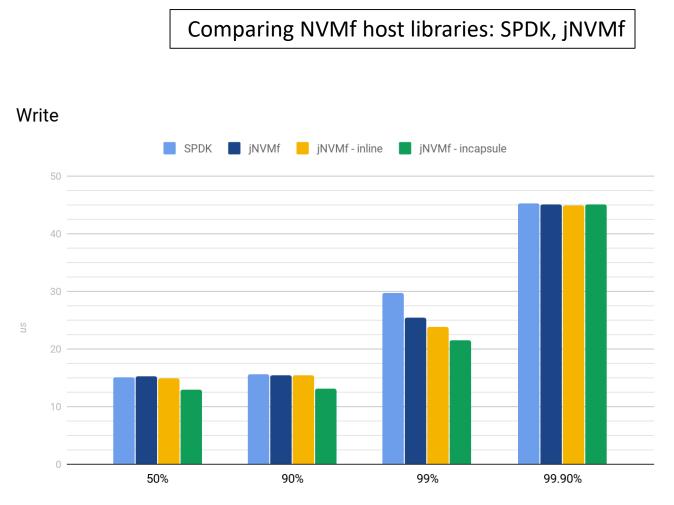
- Intel 900p device
- SPDK host
- SPDK target fastest
- Linux kernel and CX-5 offload similar delay
- Zero CPU load for CX-5 target (CPU load not shown)



### SPDK vs jNVMf Latency 4k (Write)

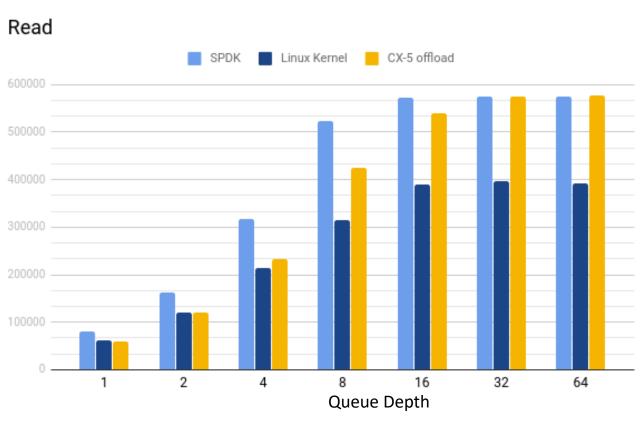
- SPDK NVMf target
- Intel 900p device
- SPDK client and jNVMf show similar performance
- No 'Java penalty'
- Enabling RDMA inline Send gives marginal advantage
- Enabling Incapsule data saves 1 RTT





### NVMf IOPS 4k (Read)

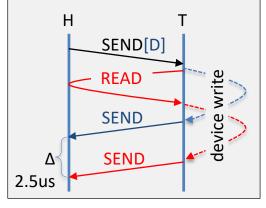
- Intel 900p device
- SPDK client
- CX-5 offload and SPDK target saturate device
- Kernel targets flattens out at ~400K IOPs
- Zero CPU load for CX-5 target (CPU load not shown)

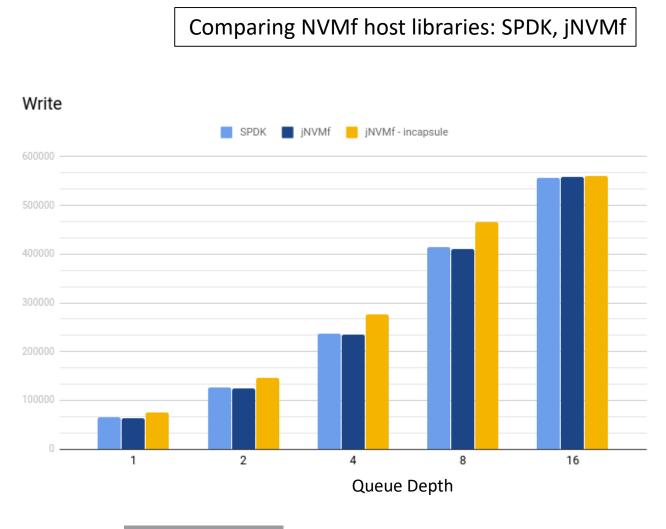


Comparing NVMf targets: SPDK, Linux Kernel, CX-5 offload

### SPDK vs jNVMf IOPS 4k (Write)

- Intel 900p device
- SPDK target
- Both SPDK and jNVMf host saturate device already at QD 16
- jNVMf incapsule data better for smaller QD





### TeraSort and TPC-DS Benchmark Setup

#### Setup

- 100Gbs RoCE: Mellanox ConnectX-5, SN2700
- 8 machines, 16 cores Intel E5-2690 @2.9Ghz
- 256GB DRAM, 128GB of it given to Spark
- NVMe devices used
  - Samsung 960Pro 1TB
- NVMf Targets deployed
  - SPDK 18.01
  - Linux Kernel 4.13
- Experiments
  - TPC-DS
  - TeraSort

#### Vanilla Spark setup:

- All data **always** kept in **DRAM**
- Input/Output to /tmpfs HDFS mount
- All intermediate operations in DRAM
- No disk spill

#### Crail NVMf setup:

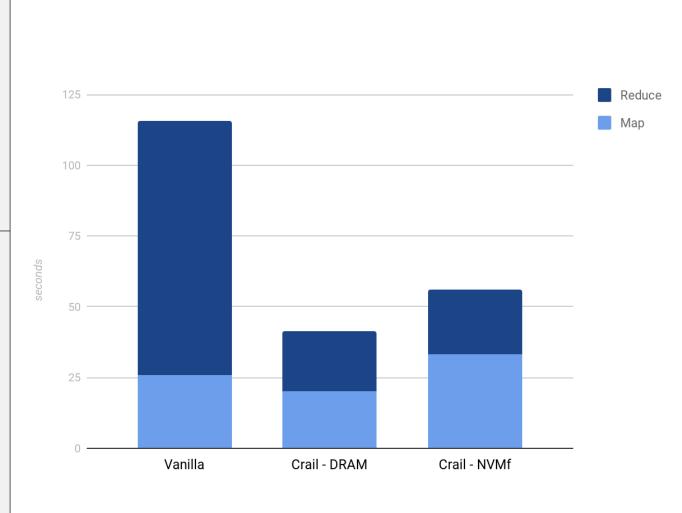
- All data read and written within NVMf
- No shuffle etc. in DRAM

### Spark TeraSort Benchmark (Spark@DRAM vs Crail@NVMf)

- Samsung 960Pro
- jNVMf client
- SPDK target
- Sorting 200GB

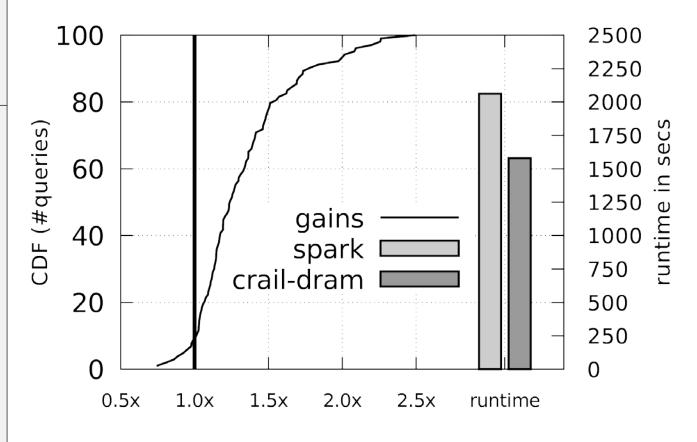
#### Spark with input/output and map/reduce in DRAM

- Crail input/output + map/reduce in NVMf
- Crail DRAM 3 times faster than vanilla Spark
- Crail 100% NVMf tier clearly outperforms vanilla Spark
- Reduce time same for Crail DRAM or NVMf
- Some penalty during Map phase (write into NVMf)



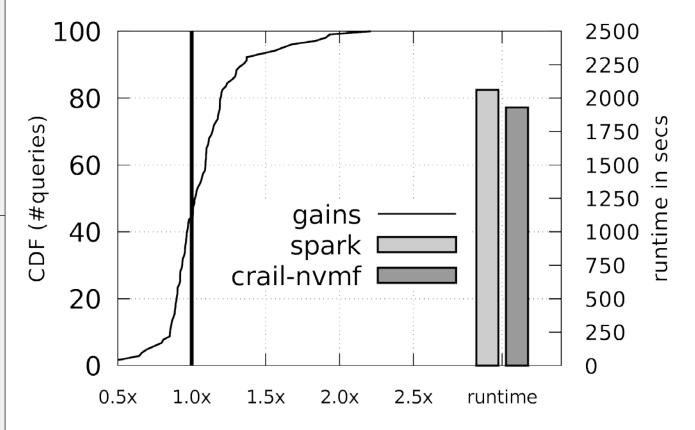
### TPC-DS Query Performance (Spark@DRAM vs Crail@DRAM)

- Spark with input/output + shuffle in DRAM
- Crail input/output + shuffle in DRAM
- Almost all queries are faster, up to 2.5x
- Crail makes more efficient use of available hardware resources



### TPC-DS Query Performance (Spark@DRAM vs Crail@NVMf)

- Samsung 960Pro devices
- jNVMf client
- Kernel NVMf target
- Spark with input/output + shuffle in DRAM
- Crail input/output + shuffle on NVMf
- Half of the queries are faster
- Overall Crail with NVMf faster than Vanilla Spark in DRAM: save cost and speed up!



### Summary & Outlook

- NVMf is the adequate extension of NVMe in a distributed system
  - Allows efficient management of ephemeral data which do not fit into DRAM
  - Using Crail + NVMf: lower cost and better performance
- NVMf supports clever device I/O management
  - In-Capsule data accelerates short NVM writes
  - Dataset management allows to pass hints to device unfortunately not yet supported in all NVMf implementations
  - Adding Bit Bucket semantic to NVMf would help, even if target device is not capable of it
- We today mainly looked at DRAM replacement
  - NVM performance is one aspect persistency is another
  - Working on NVMf tier data recovery feature (logging, replay)
- NVM tier access patterns need further investigation
  - At file level data are written sequentially, but
  - Writing huge amounts of data in parallel globally random access at device level
- Apache Crail including NVMf tier open source at https://github.com/apache/incubator-crail



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**THANK YOU** 

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