

14th ANNUAL WORKSHOP 2018

T10-DIF OFFLOAD

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IT IS ALL ABOUT INTEGRITY

Storage system is typically multi-layered

- Memory
- OS file system
- Remote storage adaptor (HBA/NIC/..)
- Storage Controller (SATA/SAS/NVME/..)

Any of the above has it's own integrity check

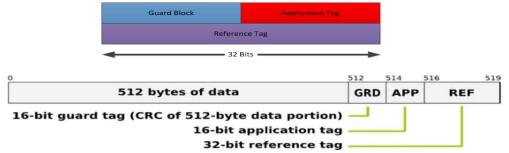
- CSUM/CRC
- Parity check
- But none protect against:
 - OS bugs, driver bugs, disk controllers and their FW errors and storage admin errors
 - Data at rest
- End to end storage integrity protection is needed

Data integrity metadata should be generated together with data creation



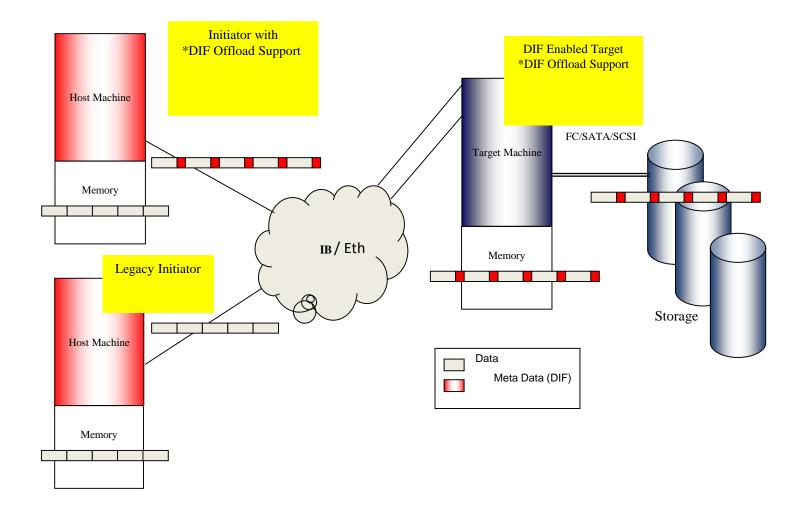
WHAT IS DIF

- DIF = Data Integrity Field
 - Aka "PI" Protection Information
 - Aka "Signature"
- The T10 standard committee specify an additional 8 byte field designated for data integrity/protection for each data block (usually of size 512 bytes but not a must).

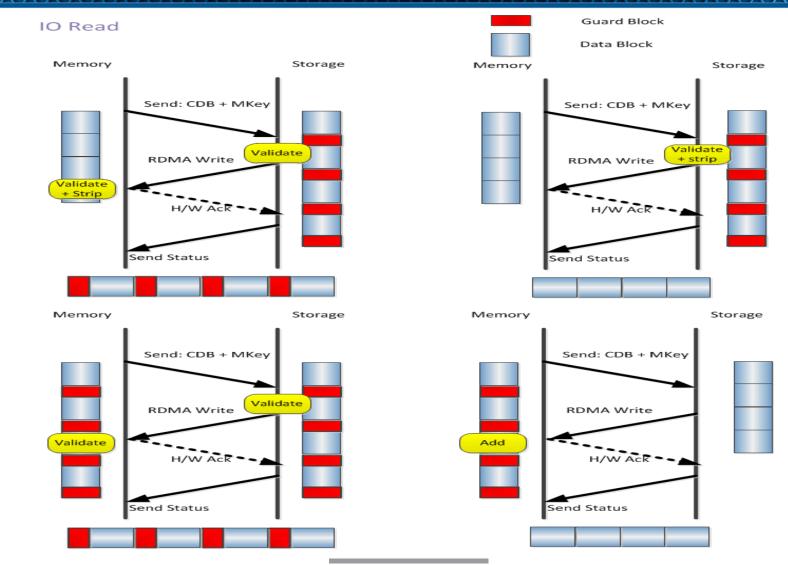


- GUARD tag (Logical Block Guarding)
 - 2 bytes CRC covering the 512 byte data sector
- APPLICATION tag (Up for grabs)
 - 2 bytes per sector ownership negotiated with target
- REFERENCE tag (Misdirected writes)
 - 4 bytes Information associated with a specific data block, typically lower 4 bytes of the Logical Block Address

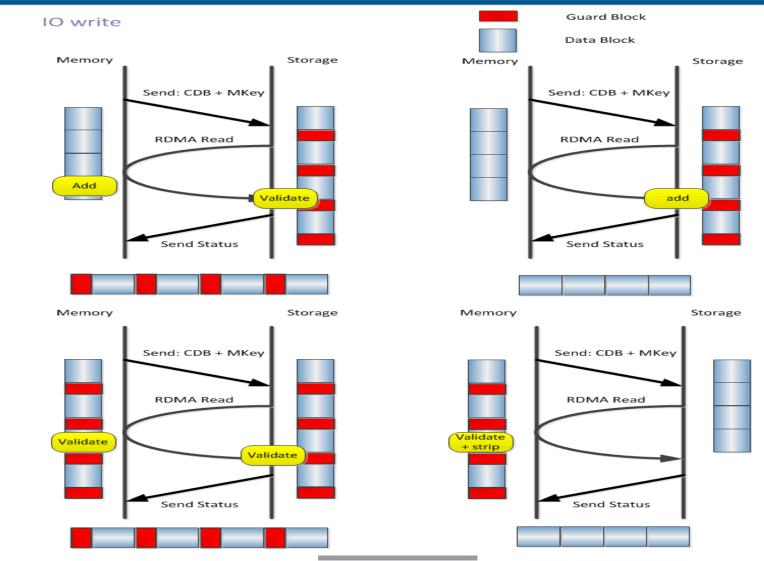
TYPICAL USE CASE



IO READ



IO WRITE



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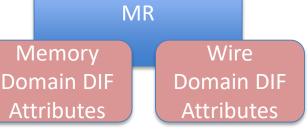
BASICS

T10-DIF is a property of MR (SIGMR)

- DIF calculations are needed when transferring data between wire and memory
- Signature is property of the memory layout
- Typically a single MR describes a single IO transaction
- Well defined for passive read/write side
- Allows single QP to serve protected and non protected transactions

Additional attributes to SIGMR describe

- How data is organized in memory
 - With/without DIF
 - DIF field is on separate buffer than data or interleaved with data
 - Block size, CRC/checksum, etc...
- How is data organized in wire
 - With/without DIF
 - Must be interleaved
 - Block size, CRC/checksum, etc...



Performing IO operation using SIGMR will always go through Signature processing

- According to SIGMR attributes
- No matter if access is local or remote
- HW will add/strip/pass the DIF info and verify it if possible

T10-DIF ERROR

DIF error is NOT a transport error

- Data transfer will not fail and may complete successfully
- Same QP may service different entities
 - QP can't transfer to error state on Signature check error

Storage app should actively check for DIF error

- Inspect the SIGMR after transaction is finished
- A lightweight operation (no device access)

Storage app should KNOW that transaction is finished

- Before checking DIF errors
- Otherwise DIF error may happen later

IBV_SET_LAYOUT_SIGNATURE

sig_attrs

Memory and wire DIF attributes

int ibv_set_layout_signature(

- struct ibv_mr *mr, int flags, struct ibv_signature_attrs sig_attrs,
- struct ibv_sge *data, struct ibv_sge *sig);

data

- A single SGE that points to the data buffer
- sge.mr can result from previous set_layout_*() calls
 - Hence can describe sophisticated memory layouts

sig

- Needed in case DIF is not interleaved with data buffers
- A single SGE that points to the data buffer
- sge.mr can result from previous set_layout_*() calls
 - Hence can describe sophisticated layouts

SIGNATURE ATTRIBUTES

- Two sets of DIF parameters
 - Memory
 - Wire
- Currently one type of DIF: T10-DIF

T10-DIF parameters:

- bg_type: CRC or IPCHECKSUM
- pi_interval: sector (block) size
- bg: block guard seed
- app_tag: application tag
- ref_tag: first block LBA
- check_mask: what to check, bit-per-byte of the T10-DIF 8 bytes
- apptag_check_mask: what to check in apptag, bit-per-bit

```
struct ibv_signature_attrs {
    struct ibv_signature_domain mem;
    struct ibv_signature_domain wire;
};
```

```
struct ibv_signature_domain {
    enum ibv_signature_type sig_type;
    union {
        struct ibv_t10dif_domain t10dif;
    };
```

```
};
struct ibv_t10dif_domain {
          enum ibv t10dif bg type bg type;
         enum ibv t10dif flags flags;
                     pi_interval;
         uint16 t
         uint16 t
                     bq;
         uint16_t
                     app_tag;
         uint32_t
                     ref_tag;
         uint8 t
                     check mask;
         uint16_t
                      apptag_check_mask;
```

```
};
```

IBV_CHECK_MR_STATUS()

Check if there was T10-DIF error on the Signature MR

If T10-DIF failed, get all needed info

- What has failed: GUARD, REFTAG or APPTAG
- What were the expected and actual values
- At which offset the error occurred
 - Will be block granularity...
- The mkey of the data block

int ibv_check_mr_status(
 struct ibv_mr *mr,
 u32 check_mask,
 struct ibv_mr_status *mr_status)

struct ibv_mr_status {
 u32 fail_status;
 struct ibv_sig_err sig_err;
};

struct ibv_sig_err {	
enum ibv_sig_err_type err_type;	
u32	expected;
u32	actual;
u64	<pre>sig_err_offset;</pre>
u32	key;
};	

PIPELINING ISSUE

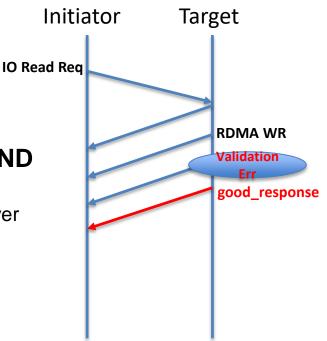
- ibv_check_mr_status() must happen at the end of transaction
- Storage systems used to respond to READ IO request by posting
 - RDMA_WRITE(data)
 - SEND(good_status)

No wait needed between RDMA_WRITE and SEND

- Because nothing can go wrong unless transport error
- Transport error kills QP, hence SEND(good_response) will never execute

• With T10-DIF, can't do that anymore

- If RDMA_WRITE generates SIG ERROR, can't stop SEND(good_response) from happening – BUG!
- Must wait for RDMA_WRITE to complete before posting SEND(good_response)



PIPELINING ISSUE RESOLVED

If SIG ERROR happens, QP revert to SQD state

 After RDMA_WRITE operation with bad signature check has finished, before SEND(good_status) started

Storage app will get notified (SQD async QP affiliated event)

Now can check MR status and see the DIF error

Two new verbs

- ibv_get_current_wrid()
 - To get the next WRE of the SEND(good_response)
- ibv_cancel_current_wr()
 - To cancel SEND(good_response) from happening

modify_qp(RTS)

- SEND(good_response) WQE was removed from the WQ
- SEND(bad_response) at a later time when convenient

SUMMARY

- End to end data integrity check is mandatory to avoid corruption
- It is way better to drop corrupted data than receive it
- Network storage systems are great use case for NIC integrity check offload
- Perfect fit for memory context attribute
- T10-DIF status and future work
 - For kernel verbs, upstream already
 - For user verbs
 - Upstreaming to rdma-core WIP
 - Older API supported since MLX_OFED 4.1 UR



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THANK YOU Tzahi Oved, Oren Duer Mellanox Technologies

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