

13th ANNUAL WORKSHOP 2017

URDMA: RDMA VERBS OVER DPDK

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EXISTING SOFTWARE RDMA DRIVERS

softiwarp and rxe

- Implement iWARP over TCP and RoCEv2, respectively
- Data transfer in kernel space
- Run unmodified verbs applications
- Designed with performance in mind

libfabric sockets provider

- Implements private protocol
- Userspace implementation using TCP/IP sockets
- Cannot run verbs applications
- High performance explicitly not a goal

SOFTWARE VERBS DRIVERS: KERNEL VS. USER SPACE

Why not implement a verbs driver using sockets API from userspace?

Userspace verbs API design choices

- Verbs will not load a userspace driver without a corresponding uverbs device exposed by the kernel
- Connection management deferred to kernel by librdmacm
- CQ events delivered from kernel
- Using userspace sockets API requires both userspace and kernel involvement

Using kernel sockets API

- Incoming RDMA READ and RDMA WRITE can be handled entirely in kernel without waking user thread
- Can use tricks like sendpage() to send TCP segments with zero-copy

Path of least resistance has been implementation using sockets API in kernel

URDMA: USERSPACE RDMA

Goals

- Prototype software RDMA driver with data transfer entirely in userspace
- Run unmodified verbs applications
- High performance

Why a userspace implementation?

- Ease of development, makes it easy to use as a development vehicle for new RDMA features
- Avoid context switches between kernel and userspace (especially for small SENDs)

Implementation uses DPDK (Data Plane Development Kit)

BACKGROUND: DPDK (DATA PLANE DEVELOPMENT KIT)

DPDK leverages Linux UIO/VFIO to map Ethernet NICs into userspace

Features:

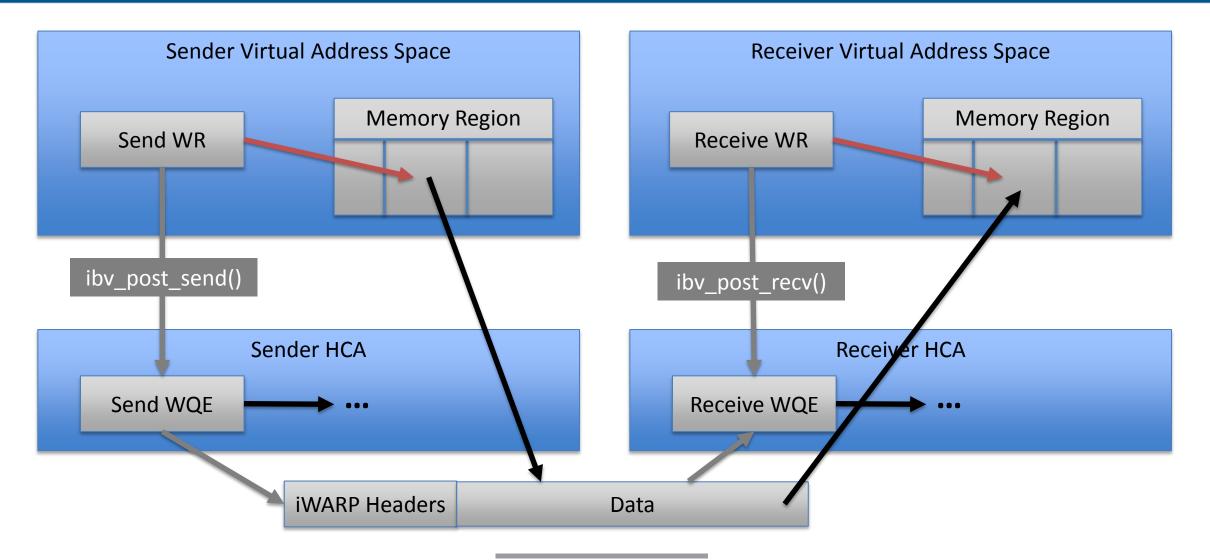
- Bulk packet transmit/receive to/from hardware NIC queues
- NUMA-aware memory buffer pool allocation using hugepages
- High performance multi-core data structures
- Hardware packet filtering
- TCP/UDP offloads, including checksum calculation

Does not provide:

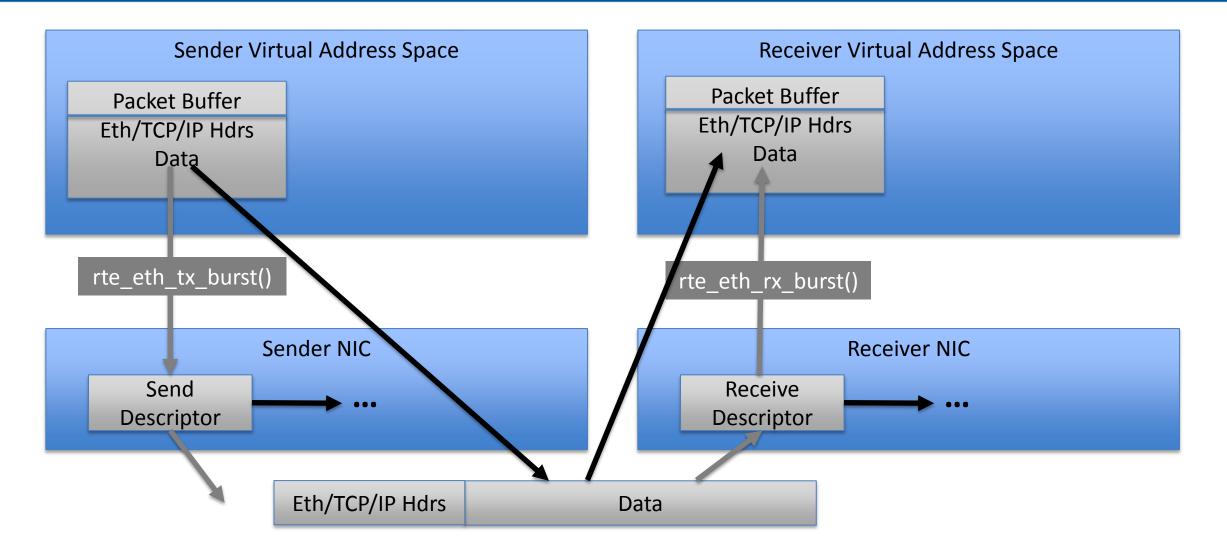
- RDMA functionality
- Network-layer or transport-layer protocol logic

Using DPDK for userspace RDMA verbs eliminates kernel from data transfer path

RDMA SEND/RECV MESSAGE TRANSFER



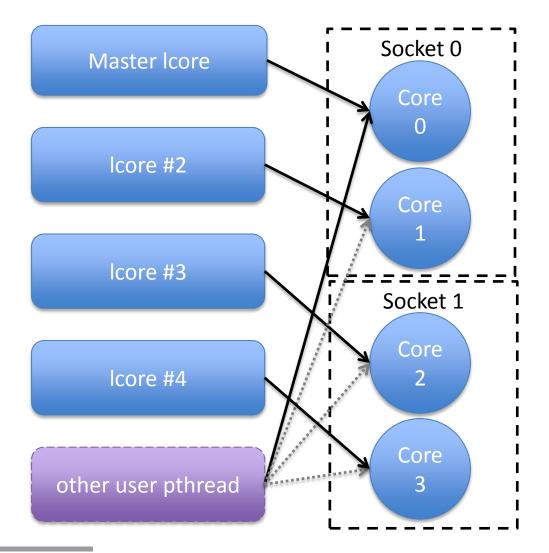
DPDK PACKET TRANSFER



BACKGROUND: DPDK THREAD MODEL

DPDK process consists of threads called "logical cores" or "lcores"

- DPDK creates 1 "lcore" thread per CPU core by default
- Thread which initializes DPDK is "master" lcore
- CPU affinity of each thread, including master, is set to run on a specific CPU core
- API allows launching tasks on other logical cores
- DPDK API expected to be called from lcores, in particular ring queues and memory pools rely on this
- We tell DPDK not to create lcores other than the master lcore

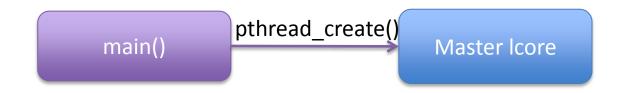


BACKGROUND: DPDK THREADS AND LIBRARIES

DPDK is more of an application framework than a library

DPDK initialization function

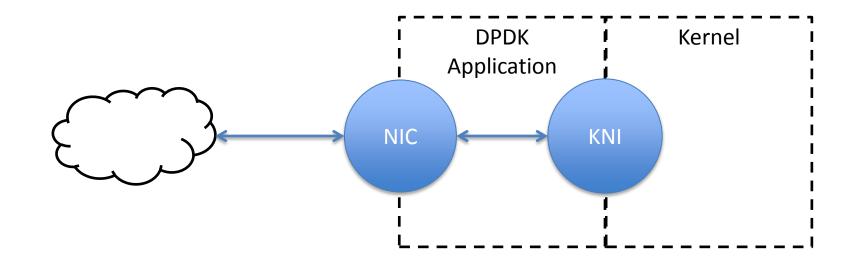
- Takes command-line arguments
- Consumes all available hugepages by default
- Changes CPU affinity of calling thread
- To use DPDK from library, we create a thread and call DPDK initialization from there
 - Pass parameter to not create further lcores
 - Separate DPDK thread from user threads
 - We do not affect CPU affinity of user threads



BACKGROUND: DPDK KNI

KNI (Kernel Network Interface)

- Creates a virtual network interface in the kernel
- Loosely associated with a DPDK Ethernet hardware NIC
- Can exchange packets between kernel and userspace
- Useful for small interactions between kernel service and DPDK application





URDMA: DESIGN AND IMPLEMENTATION

URDMA: DESIGN

Implements iWARP DDP and RDMAP protocols

Runs over UDP transport protocol

- TRP (Trivial Reliability Protocol) provides a thin shim for reliability
- Simplifies implementation considerably

Small kernel component

- Required for libibverbs initialization, RDMA CM, and CQ events
- Performs connection establishment before ceding control of UDP "connection" to liburdma
- Uses KNI to send/receive packets to/from userspace

Packet processing done in background thread

• Ensure quick response to RDMA packets and KNI events

Hardware receive filter used to assign queue pairs to NIC receive queues

URDMA: MULTI-PROCESS SUPPORT

• DPDK maps Ethernet NIC hardware into userspace \rightarrow owned by that process

- Can delegate to secondary processes that explicitly cooperate
- DPDK considers primary + secondary processes as one combined application
- DPDK threads in combined application cannot share the same lcore identifier

In urdma, primary process is a user daemon urdmad

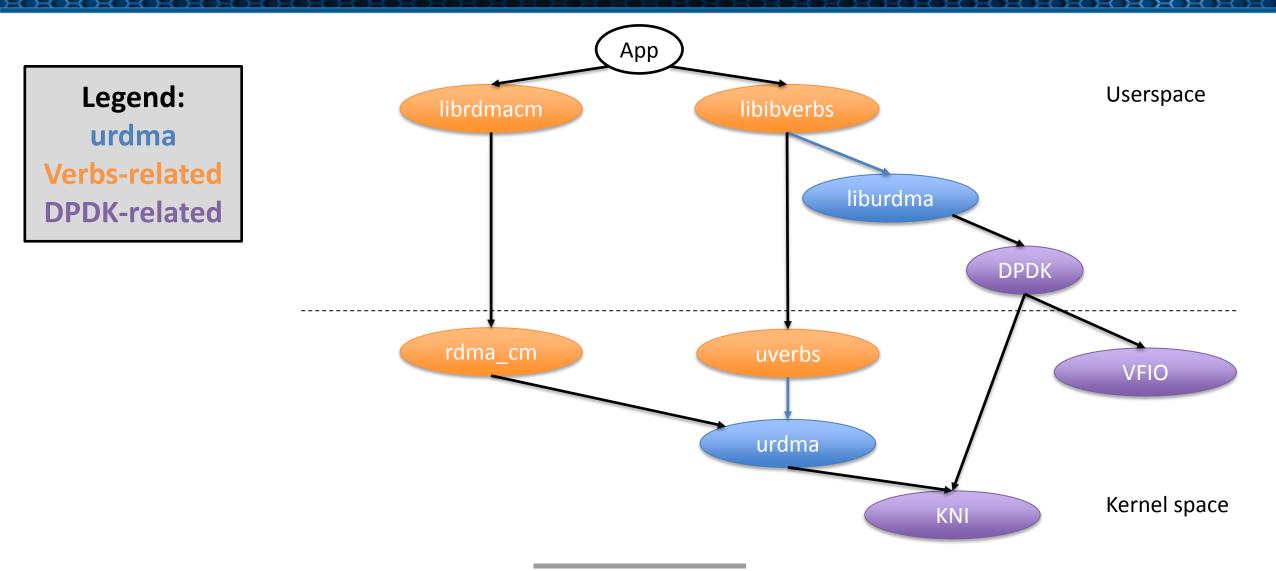
- Initializes DPDK
- Registers secondary processes with separate core mask
- Assigns Ethernet NIC hardware RX/TX queues to urdma processes
- Sets up Ethernet NIC hardware filtering rules

liburdma verbs provider

- Sets up process as secondary DPDK process
- DPDK "master" lcore acts as background progress thread

Each liburdma process has direct access to its Ethernet NIC hardware queues

URDMA: COMPONENTS



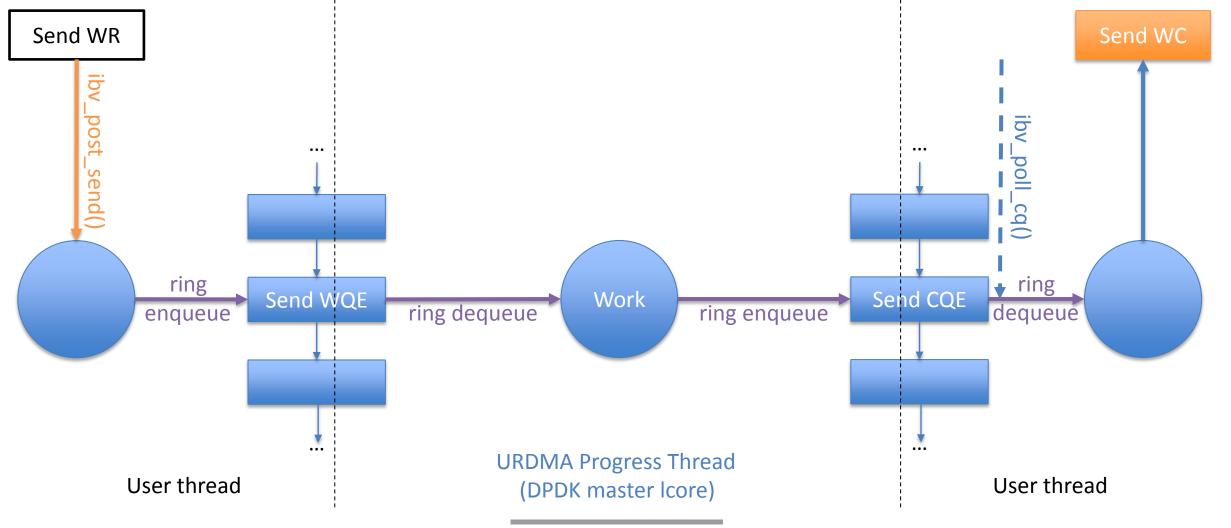
Connection establishment done in kernel space

In userspace:

- Each queue pair must be assigned a Ethernet NIC hardware send and receive queue
- Hardware receive filtering rules must be assigned before first packet arrives
- Private character device used to communicate connection establishment

арр		ma	urdmad	kernel	rdma_cn uverbs
	Create QP	Assign TX/RX queue			urdma
		ibv_cmd_create_qp()		>	
R	DMA Connect/Accept				
				QP connected	
				QP Ready to recv	
			RDMA CM	1 Established Event	

URDMA DATA TRANSFER







PERFORMANCE: OVERVIEW

• Two identical systems:

- Supermicro SYS-6028R-T
- 2 Intel Xeon ES-2630 v4 CPU @ 2.20GHz
- 64 GB DDR4 RAM
- PCIe generation 3
- Ubuntu 16.10 with inbox 4.8 kernel
- Intel XL710 40GbE NIC
- Verbs and RDMA CM as supplied with Ubuntu 16.10

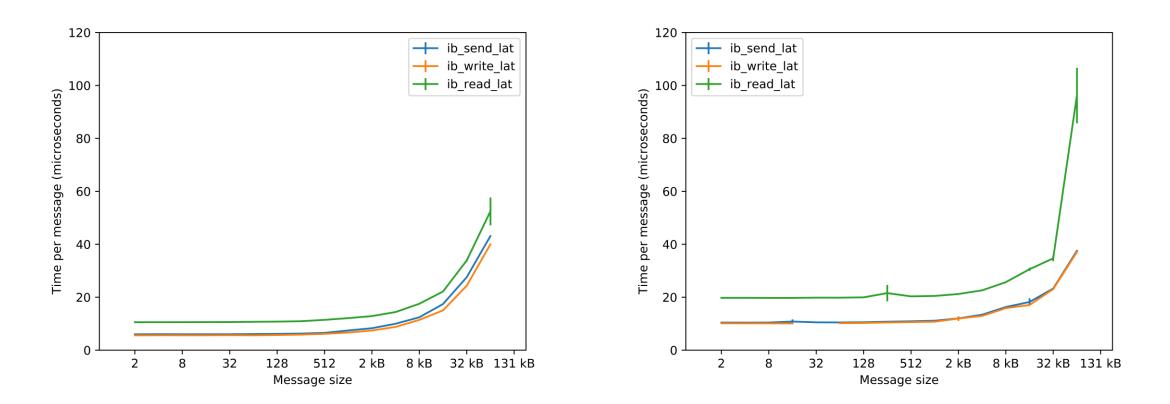
Applications used

- perftest version 3.0+0.18.gb464d59-1
- UNH EXS (Extended Sockets) 1.4.1 (<u>https://www.iol.unh.edu/expertise/unh-exs</u>)

RAW VERBS: LATENCY

urdma

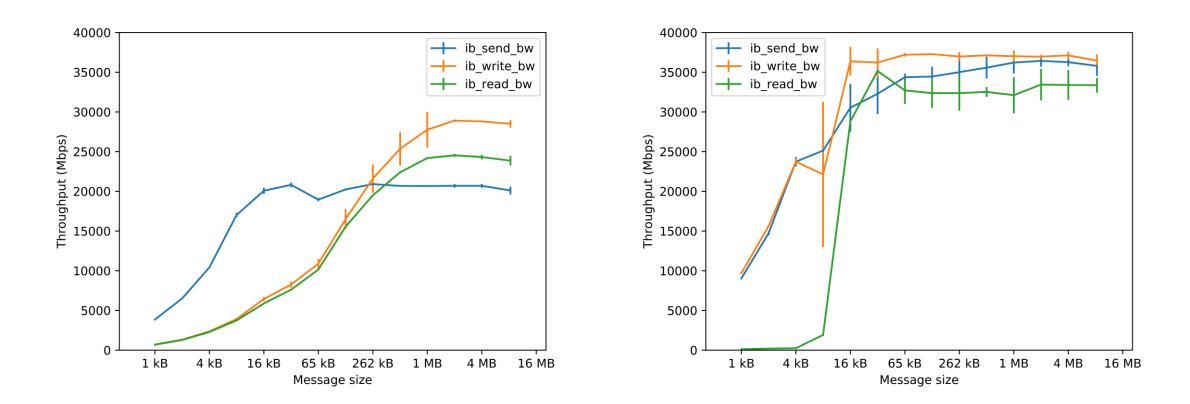




RAW VERBS: THROUGHPUT

urdma

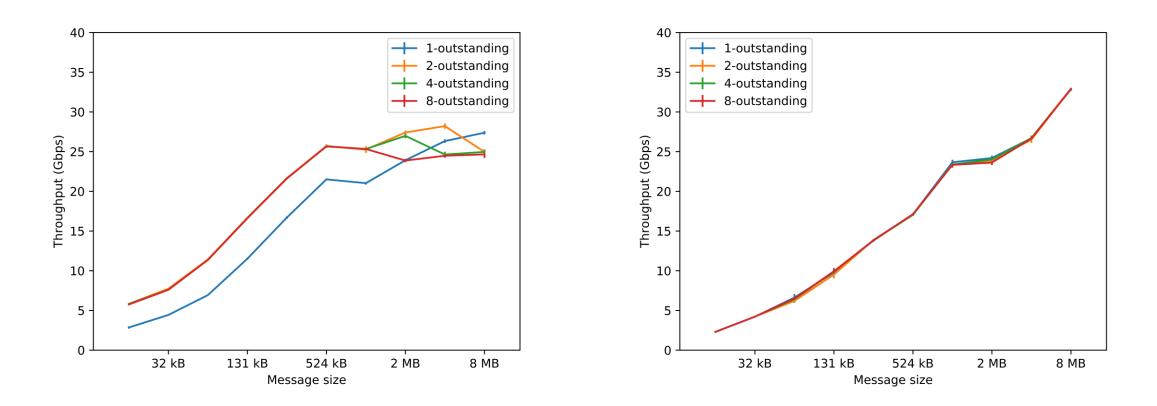
softiwarp



UNH EXS: THROUGHPUT

urdma

softiwarp







URDMA SUMMARY

- Existing software RDMA implementations done in kernel space
- DPDK allows us to implement RDMA verbs data transfer in userspace
 - Eliminates all kernel involvement in data transfer path
 - Small kernel module for connection management
- Runs unmodified verbs applications
- Designed with performance in mind
 - Good small message latency
 - Needs tuning for throughput

Future work

- Investigate zero-copy transmit support
- libfabric provider implementation
- Reliable datagram support

URDMA DOWNLOAD AND STATUS

- urdma development done on GitHub
 - <u>https://github.com/zrlio/urdma</u>
- No formal release as of yet
- Not integrated into rdma-core
- Tested on Ubuntu 16.10 and DPDK 16.07



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THANK YOU

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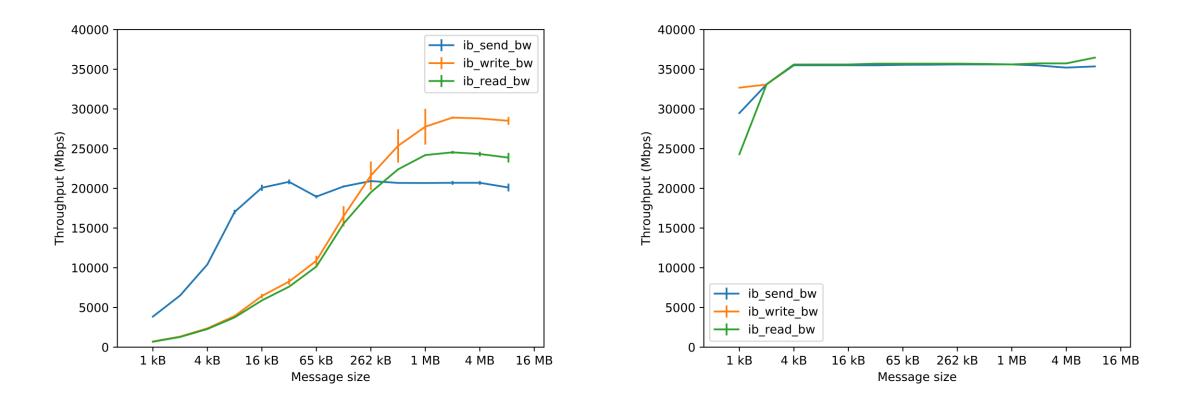




RAW VERBS: THROUGHPUT VS. HARDWARE RNIC

urdma on Intel XL710

Chelsio T580-LP-CR iWARP



RAW VERBS: LATENCY VS. HARDWARE RNIC

urdma on Intel XL710

Chelsio T580-LP-CR iWARP

